

VERTEX COVER MEDIANTE ARROTONDAMENTO

PROBLEMA LP

INPUT: $A \in \mathbb{Q}^{m \times n}$, $\underline{b} \in \mathbb{Q}^m$, $\underline{c} \in \mathbb{Q}^n$

SOL. AMMISSIBILI: $\underline{x} \in \mathbb{Q}^n$

$$\text{t.c. } A\underline{x} \geq \underline{b}$$

FUNZ. OBIETTIVO: $\underline{c}^T \underline{x}$

TIPO: MIN

PROBLEMA ILP

INPUT: $A \in \mathbb{Q}^{m \times n}$, $\underline{b} \in \mathbb{Q}^m$, $\underline{c} \in \mathbb{Q}^n$

SOL. AMMISSIBILI: $\underline{x} \in \mathbb{Z}^n$

$$\text{t.c. } A\underline{x} \geq \underline{b}$$

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TIPO: MIN

[LP]

- Risolvibile in tempo polinomiale in

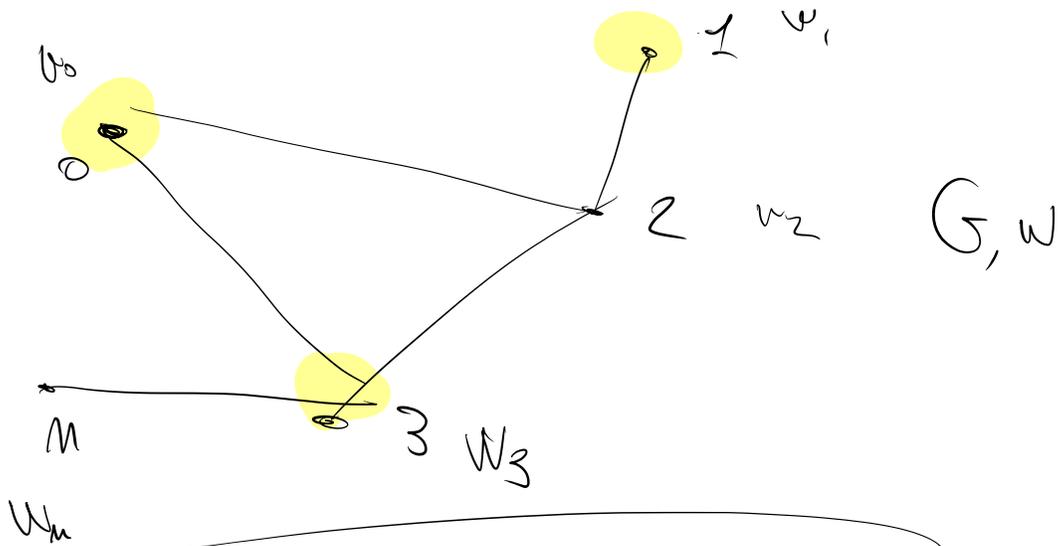
(n) [1981, 1984] [metodo del
punto interno di Karumkar]

- Dantzig

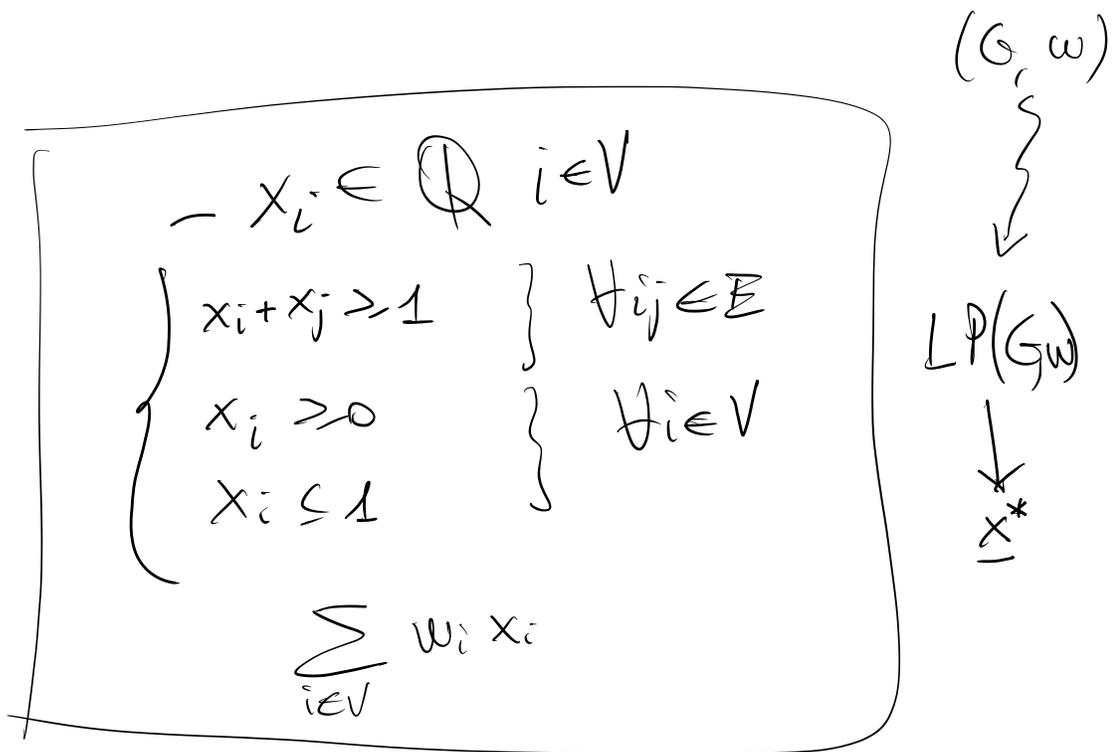
LP EPO

[ILP]

ENPO - completo



$$\begin{array}{l}
 \text{ILP}(G, w) \\
 \left. \begin{array}{l}
 - x_i \in \mathbb{Z} \quad i \in V \\
 \left. \begin{array}{l}
 x_i + x_j \geq 1 \\
 x_i \geq 0 \\
 x_i \leq 1
 \end{array} \right\} \forall ij \in E \\
 \left. \begin{array}{l}
 \\
 \\
 \end{array} \right\} \forall i \in V
 \end{array} \right\} \\
 \sum_{i \in V} w_i x_i
 \end{array}$$



Lemma 1: $w_{LP}^* \leq w_{ILP}^*$

$\pi_i = \begin{cases} 1 \\ 0 \end{cases}$ $x_i^* \geq \frac{1}{2}$

$\geq \text{Anzahl}$

Lemma 2: π ist eine zulässige Lösung des ILP.

Denn: $\forall ij \in E \quad \pi_i + \pi_j \geq 1$

Per sberdo $\exists j \in E$ t.c. $r_i + r_j < 1$.

$$\Rightarrow r_i = r_j = 0$$

$$\Rightarrow x_i^* < \frac{1}{2} \text{ e } x_j^* < \frac{1}{2}$$

$$\Rightarrow \underline{\underline{x_i^* + x_j^* < 1}}$$

imposs.



$$r_i \leq 2x_i^* \begin{cases} r_i = 0 & \text{ovvia} \\ r_i = 1 \Rightarrow x_i^* \geq \frac{1}{2} \\ & 2x_i^* \geq 1 = r_i \end{cases}$$

Lemma 2:

$$\sum_{i \in V} w_i r_i \leq 2 \sum_{i \in V} w_i x_i^* =$$

$$= 2 w_{LP}^* \stackrel{\text{Lemma 1}}{\leq} 2 w_{ILP}^*$$