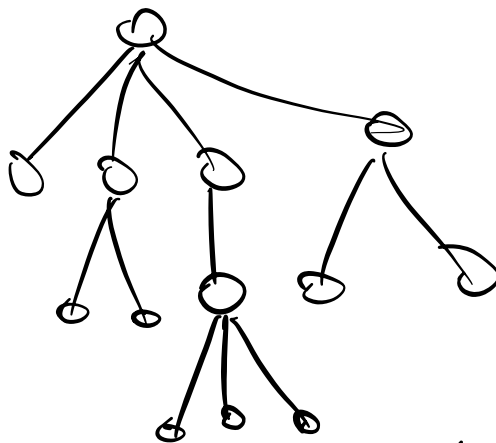


STRUTTURE SUCCINTE PER ALBERI BINARI

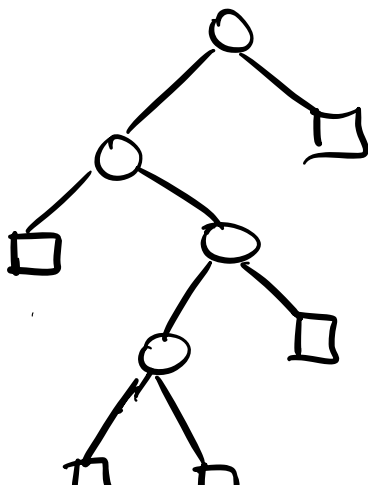
- Teoria dei grafi: albero =
grafo non orientato
senza ciclo
- Informatica



ALBERI
RADICATI
E
ORDINATI

ALBERI
GENERALI

- ALBERI BINARI



○ nodi interni
□ nodi esterni
(foglie)

$$\text{Thm: } \# \text{ nodi esterni} = \# \text{ nodi interni} + 1$$

$$|E| = |I| + 1$$

Thm: Numero di alberi con n nodi interni è

$$C_n = \frac{1}{n+1} \binom{2n}{n} \left[\begin{array}{l} \text{NUMERO DI} \\ \text{CATRANO} \end{array} \right]$$

Approssimazione di Stirling:

$$x! \sim \sqrt{2\pi x} \left(\frac{x}{e}\right)^x$$

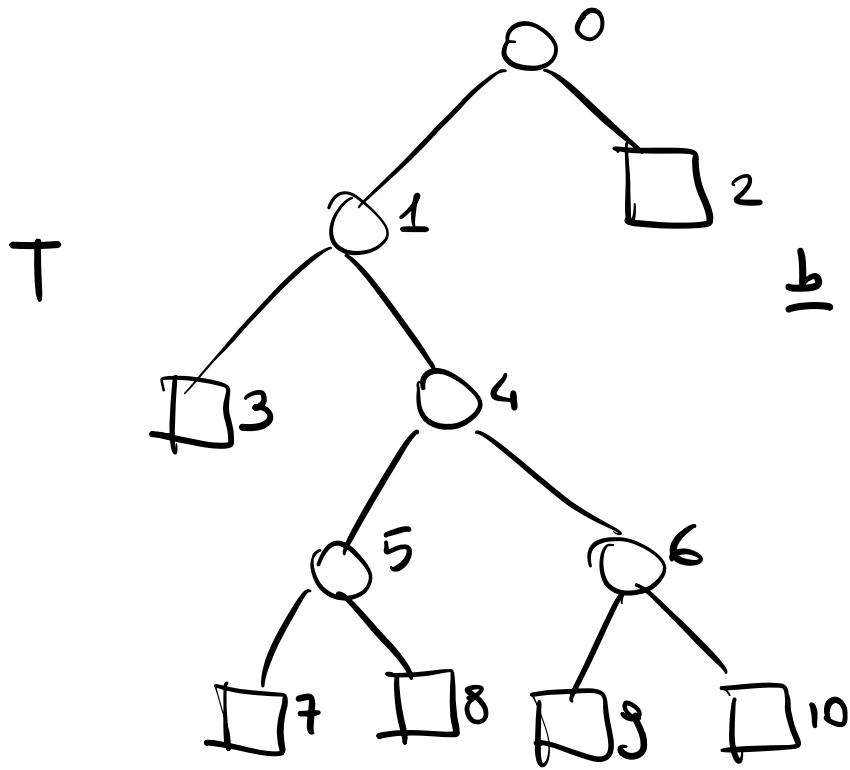
$$C_n = \frac{1}{n+1} \binom{2n}{n} = \frac{1}{n+1} \frac{(2n)!}{n! n!} \approx$$

$$\approx \frac{1}{n+1} \frac{\sqrt{4\pi n} \left(\frac{2n}{e}\right)^{2n}}{2\pi n \left(\frac{n}{e}\right)^{2n}} =$$

$$= \frac{1}{n+1} \frac{1}{\sqrt{\pi n}} 2^{2n} \times \frac{4^n}{\sqrt{\pi n^3}}$$

$$\lg C_n \approx \boxed{2n} - O(\lg n)$$

RAPPRESENTAZIONE
SUCCINTA



$$n=5$$

$$\#nodes = 2n+1 = 11$$

$$\underline{\underline{b}}$$

0	1	2	3	4	5	6	7	8	9	10
1	1	0	0	1	1	1	0	0	0	0

+ RANK/SELECT

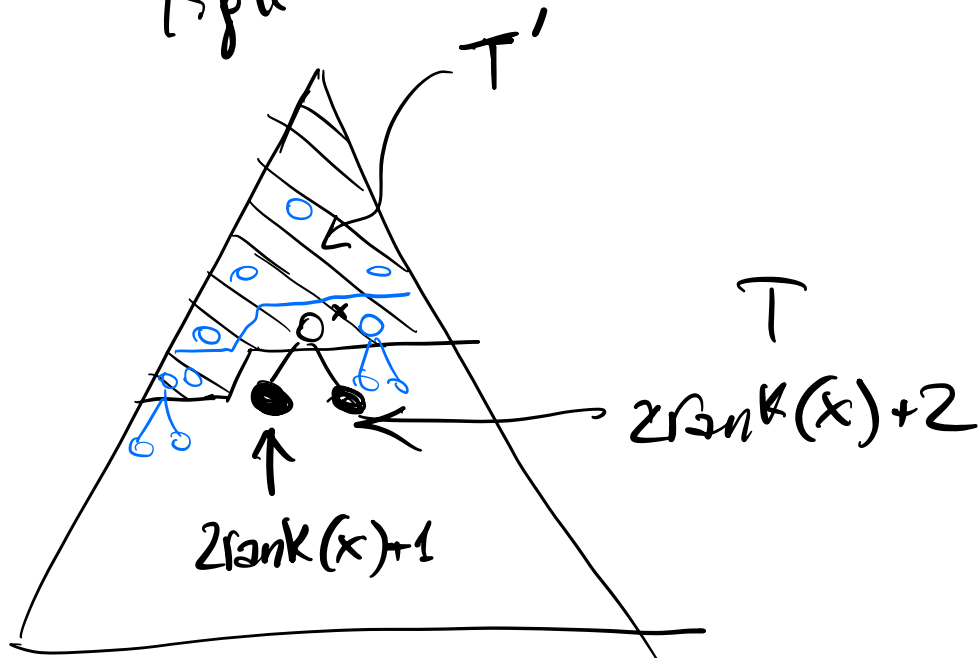
$$2n+1 + o(2n+1)$$

$$= \boxed{2n + o(n)}$$

NAVIGARE NELL'ALBERO

Se x è l'indice di un nodo
1) come capire se è
una foglia

2) ~~chi~~ sono gli indici dei figli



ACCESSO
AI
FIGLI

$$\begin{aligned}
 \# \text{ nodi interni di } T' &= \\
 &= \# \text{ nodi interni di } T \text{ con} \\
 &\quad \text{indici} < x = \\
 &= \# \text{ uni in } \underline{b} \text{ di indici} < x = \\
 &= \text{rank}_{\underline{b}}(x)
 \end{aligned}$$

$$\begin{aligned}
 \# \text{ nodi } T' &= 2 \# \text{ nodi interni di } T' + 1 \\
 &= 2 \text{rank}_{\underline{b}}(x) + 1
 \end{aligned}$$

ACCESSO AL
GENITORE

Il padre di x è un nodo
p t.c.

$$2\text{rank}(p) + 1 = x$$

opp. $2\text{rank}(p) + 2 = x$

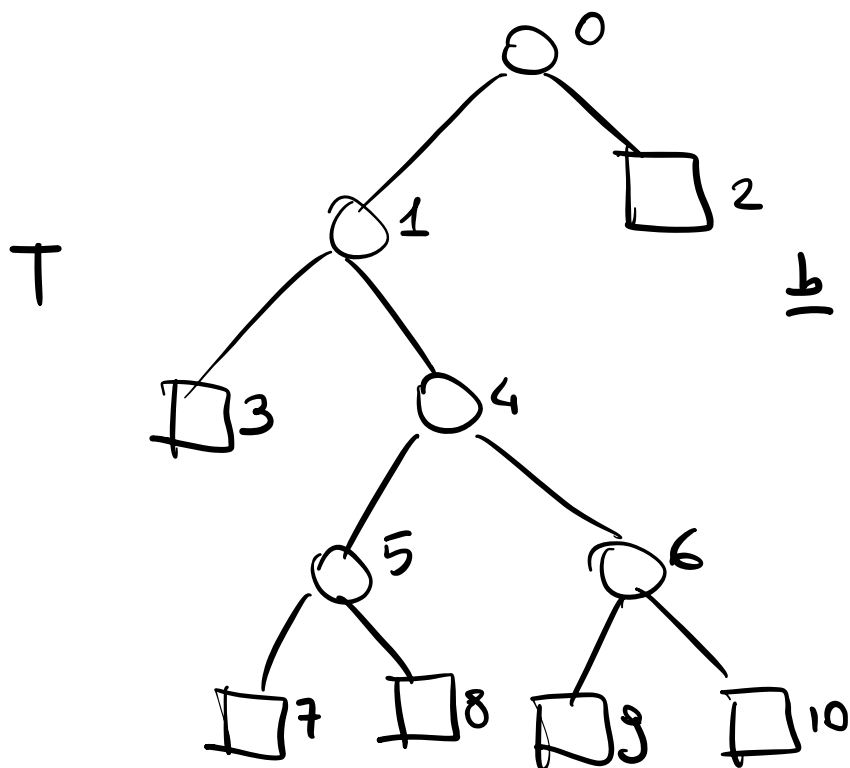
$$\left. \begin{aligned} \text{rank}(p) + \frac{1}{2} &= \frac{x}{2} \\ \text{opp. } \text{rank}(p) + 1 &= \frac{x}{2} \end{aligned} \right\}$$

$$\text{rank}(p) = \left\lfloor \frac{x}{2} - \frac{1}{2} \right\rfloor$$

$$\text{select}(\text{rank}(p)) = \text{select}\left(\left\lfloor \frac{x}{2} - \frac{1}{2} \right\rfloor\right)$$

$$p = \text{select}\left(\left\lfloor \frac{x}{2} - \frac{1}{2} \right\rfloor\right)$$

DATI ANCILLARI



$$\underline{b}$$

0	1	2	3	4	5	6	7	8	9	10
1	1	0	0	1	1	1	0	0	0	0

+ RANK/SELECT

DATI ANCILLARI ANCHE
SULLE FOGLIE

A

0	1	2				10
~	~	~	...			~

DATI ANCILLARI SOLO
NODI INTERNI

SU

B

0	1	2	3	4	5
~	~	~	~	~	~

BILIEZIONI FRA
ALBERI BINARI,
ALBERI GENERALI E
FORESTE

B_n insieme alberi
binari con
 n nodi interni

F_n insieme delle
foreste ordinate
con n nodi

T_n insieme degli
alberi generali
con n nodi

$$\varphi: F_n \xrightarrow{\cong} T_{n+1}$$

LFT



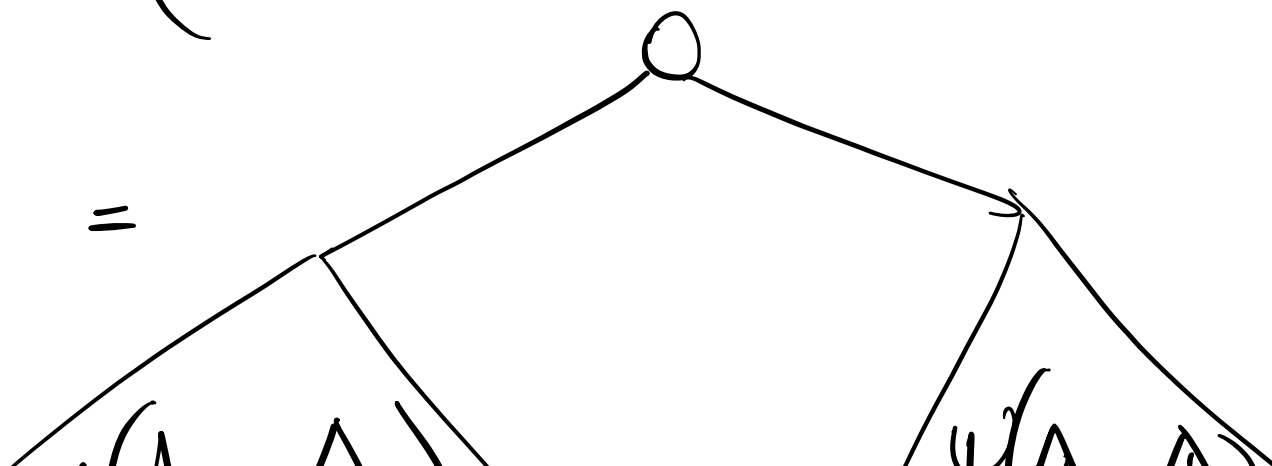
$$\psi: F_n \xrightarrow{\cong} B_n$$

FCNS (FIRST-CHILD NEXT-SIBLING)

$$\psi(\emptyset) = \square$$

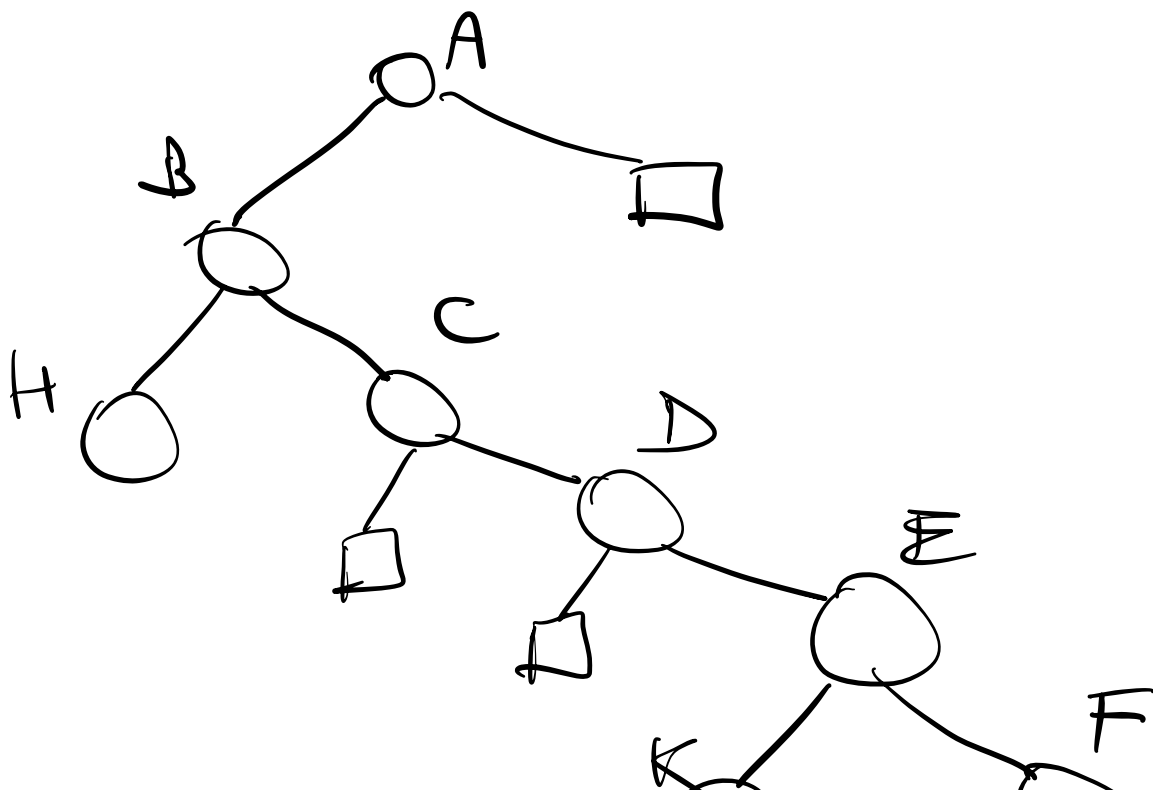
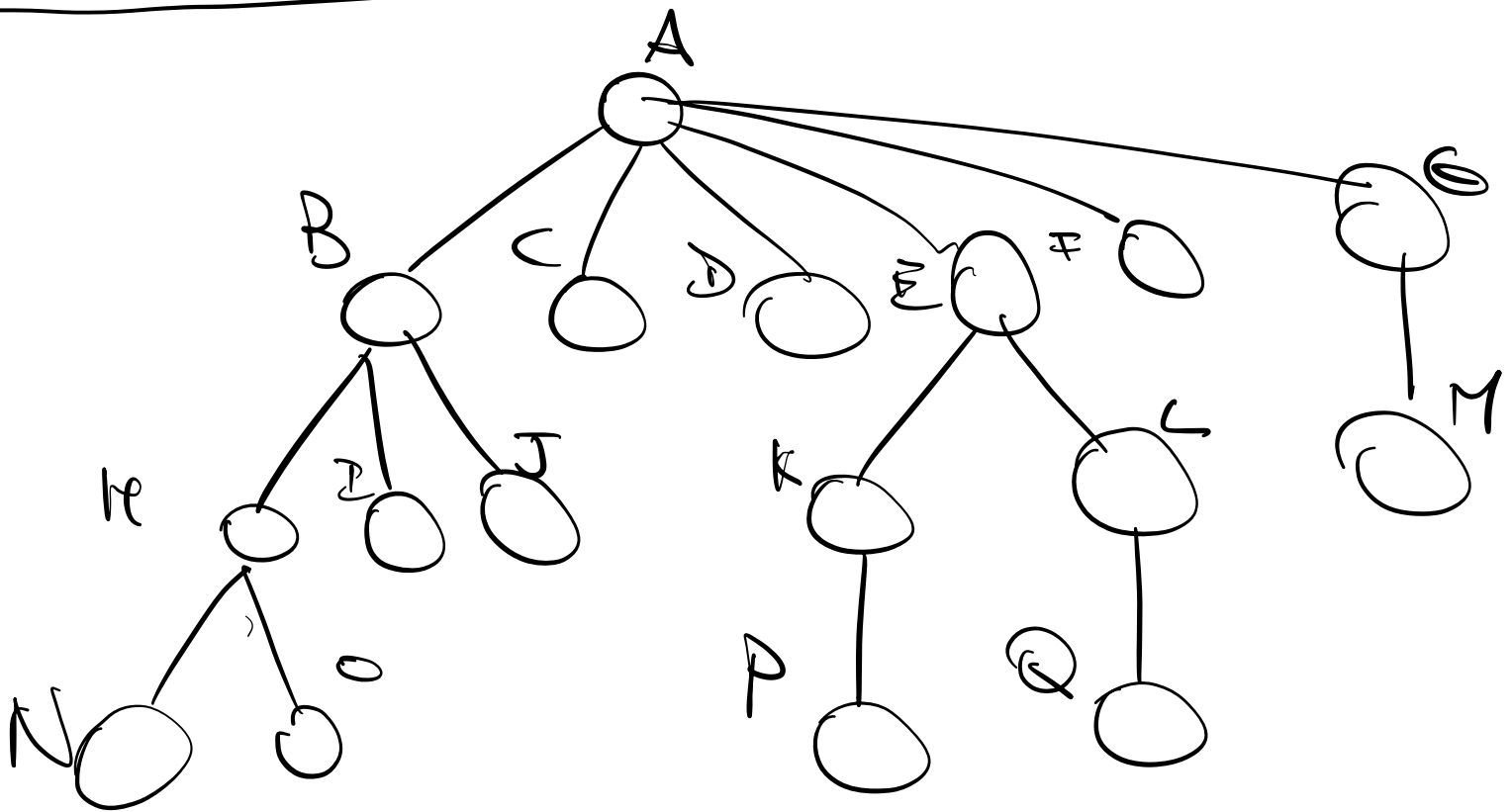
$$\psi \left(\begin{array}{c} \text{Tree with root and children } T_1, \dots, T_h \\ T'_1, \dots, T'_k \end{array} \right) =$$

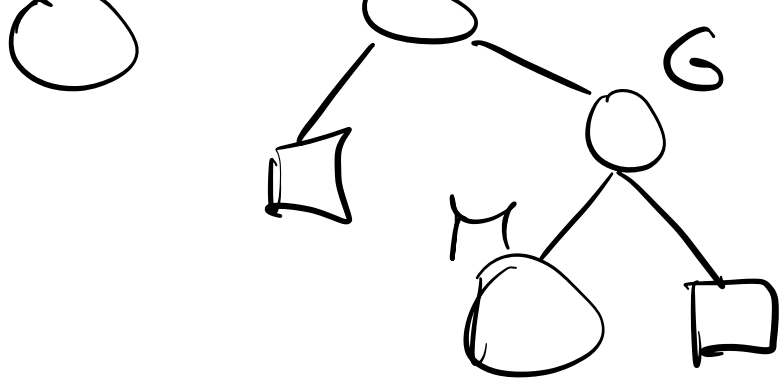
=



$$\psi(\underline{A_1} \dots \underline{A_n})$$

$$\psi(\underline{A_1} \dots \underline{A_n})$$





$$|T_{n+1}| = |F_n| = |B_n| = C_n$$

$\nwarrow$ $\nwarrow$ $\nwarrow$

PAROLE DI DYCK

$$\Sigma = \{ \underset{0}{(}, \underset{1}{)} \}$$

$w \in \Sigma^*$ è una parola di Dyck

1) $\#_{(} w = \#_{)} w$

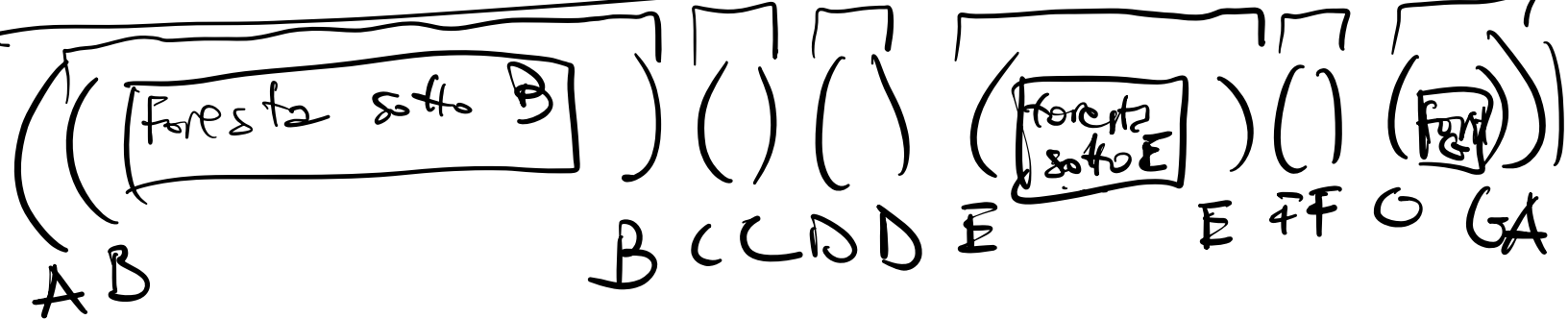
2) $\forall v$ prefisso di w

↙ n aperte

↙ n aperte

↙ n aperte





$$|D_n| = |F_n| = C_n$$

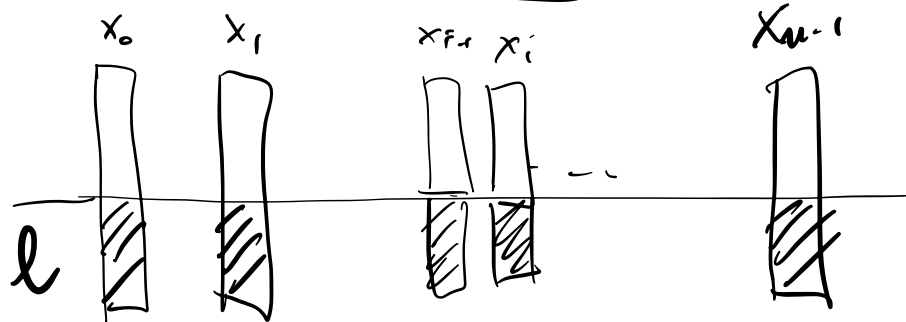
UNA RAPPRESENTAZIONE
SUCCINTA PER
SEQUENZE MONOTONE

$$0 \leq x_0 \leq x_1 \leq x_2 \leq \dots \leq x_{n-1} < U$$

$i \rightsquigarrow x_i$

RAPPRESENTAZIONE
DI ELIAS-FANO

BIT
PER SIGN.



BIT
PER SIGN.

$\log U$
RAPPRESENTAZIONE
ESPLICITAMENTE

$$l = \max \left(0, \left\lfloor \log \frac{u}{n} \right\rfloor \right)$$

$$\left[\begin{array}{l} l_0 = x_0 \bmod 2^l \\ l_1 = x_1 \bmod 2^l \\ \vdots \\ l_{n-1} = x_{n-1} \bmod 2^l \end{array} \right] \text{ LOWER PARTS}$$

$$(x_{-1} = 0)$$

$$u_i = \left\lfloor \frac{x_i}{2^l} \right\rfloor - \left\lfloor \frac{x_{i-1}}{2^l} \right\rfloor \leftarrow \text{IN UNARIO}$$

$$\text{RANK-SELECT} \leftarrow \underline{b} = \rightarrow \underbrace{0000}_{u_0} + \underbrace{0000}_{u_1} + \underbrace{0001}_{u_2} + 1$$

SPACE OCCUPATION

LOWER
UPPER

BIT
BIT

lm BIT

$$\leq \sum_{i=0}^{n-1} (u_i + 1) =$$

$$= \sum_{i=0}^{n-1} \left(\left\lfloor \frac{x_i}{2^l} \right\rfloor - \left\lfloor \frac{x_{i-1}}{2^l} \right\rfloor + 1 \right) =$$

$$\quad \quad \quad n-1 \quad (x_{n-1}) \quad (x_{i-1})$$

$$= n + \sum_{i=0}^{\infty} \left(\left\lfloor \frac{x_i}{2^i} \right\rfloor - \left\lfloor \frac{x_{i+1}}{2^i} \right\rfloor \right)$$

$$= n + \left\lfloor \frac{x_{n-1}}{2^l} \right\rfloor \leq n + \left\lfloor \frac{U}{2^l} \right\rfloor \leq$$

$$\leq n + \frac{U}{2^l} \leq n + \frac{U}{2^{\lceil \lg \frac{U}{n} \rceil}} =$$

$$\leq n + \frac{U}{2^{\lg \frac{U}{n} - 1}} = n + \frac{2U}{U/n} =$$

$$= 3n$$

$$D_n = 2n + 3n \quad \text{BIT}$$

$$= 3n + n \left\lceil \lg \frac{U}{n} \right\rceil =$$

$$= 2n + n \left\lceil \lg \frac{U}{n} \right\rceil$$

$$\text{select}_{\underline{b}}(i) = i-1 + \mu_0 + \mu_1 + \dots + \mu_i =$$

$$\mu_i = \left\lfloor \frac{x_i}{2^l} \right\rfloor - \left\lfloor \frac{x_{i-1}}{2^l} \right\rfloor \leftarrow \text{UNARIO}^{\text{IN}}$$

$$- \underline{b} = \rightarrow \underbrace{\text{oooo}}_{\mu_0} + \underbrace{\text{oooo}}_{\mu_1} + \underbrace{1 \text{ oooo}}_{\mu_2} + \dots$$

$$= i-1 + \left\lfloor \frac{x_0}{2^l} \right\rfloor + \left\lfloor \frac{x_1}{2^l} \right\rfloor - \left\lfloor \frac{x_0}{2^l} \right\rfloor +$$

$$+ \left\lfloor \frac{x_2}{2^l} \right\rfloor - \left\lfloor \frac{x_1}{2^l} \right\rfloor + \dots + \left\lfloor \frac{x_i}{2^l} \right\rfloor -$$

$$- \left\lfloor \frac{x_{i-1}}{2^l} \right\rfloor = i-1 + \left\lfloor \frac{x_i}{2^l} \right\rfloor$$

$$\text{select}_{\underline{b}}(i) = i-1 + \left\lfloor \frac{x_i}{2^l} \right\rfloor$$

$$\Rightarrow \left\lfloor \frac{x_i}{2^l} \right\rfloor = \text{select}_{\underline{b}}(i) - i + 1$$

✓ [2nd]

✓