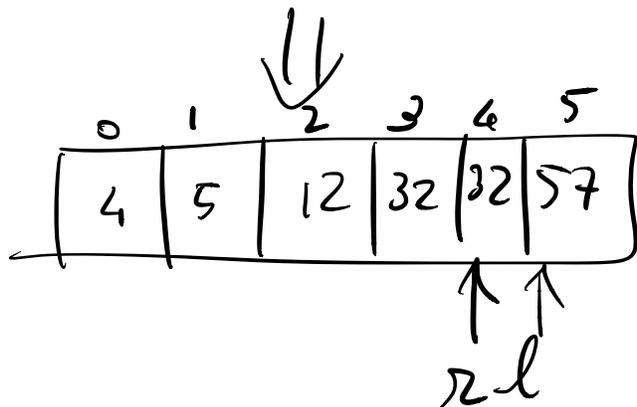
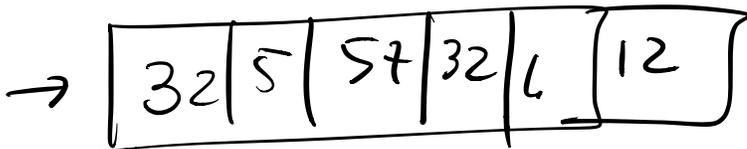
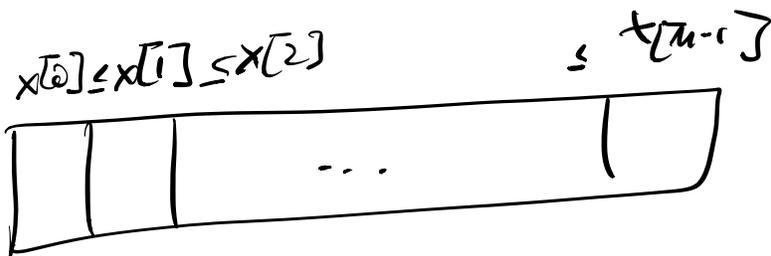


# ORDINAMENTO DI SLICE

var x [l:at



Donald Knuth  
TAOCP

36

# RICERCA

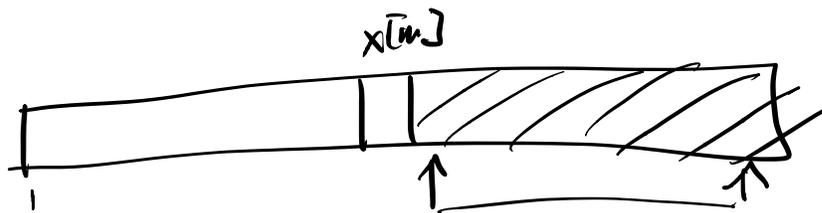
var x [ ] int  
var y int

```
func search (x [ ] int, y int) int {  
  for i, e := range x {  
    if e == y {  
      return i  
    }  
  }  
  return -1  
}
```

RICERCA  
LINEARE

slice lunga n

~~Caso PEGGIORE → n confronti~~  
~~(quando l'elemento non c'è)~~



func

binarySearch (x [int, y int) int {

var left, right int

left = 0

right = len(x) - 1

for left <= right {

m := (left + right) / 2

if x[m] == y {

return m

} else x[m] < y {

left = m + 1

} else {

right = m - 1

}

}

return -1

}

Dato  $n$ , qual è il valore  $k$   
per cui

$$\left( \frac{\binom{n}{2}}{2} \right) \left\{ k \text{ volte} < 1 \right.$$

$$\Rightarrow \frac{n}{2^k} < 1$$

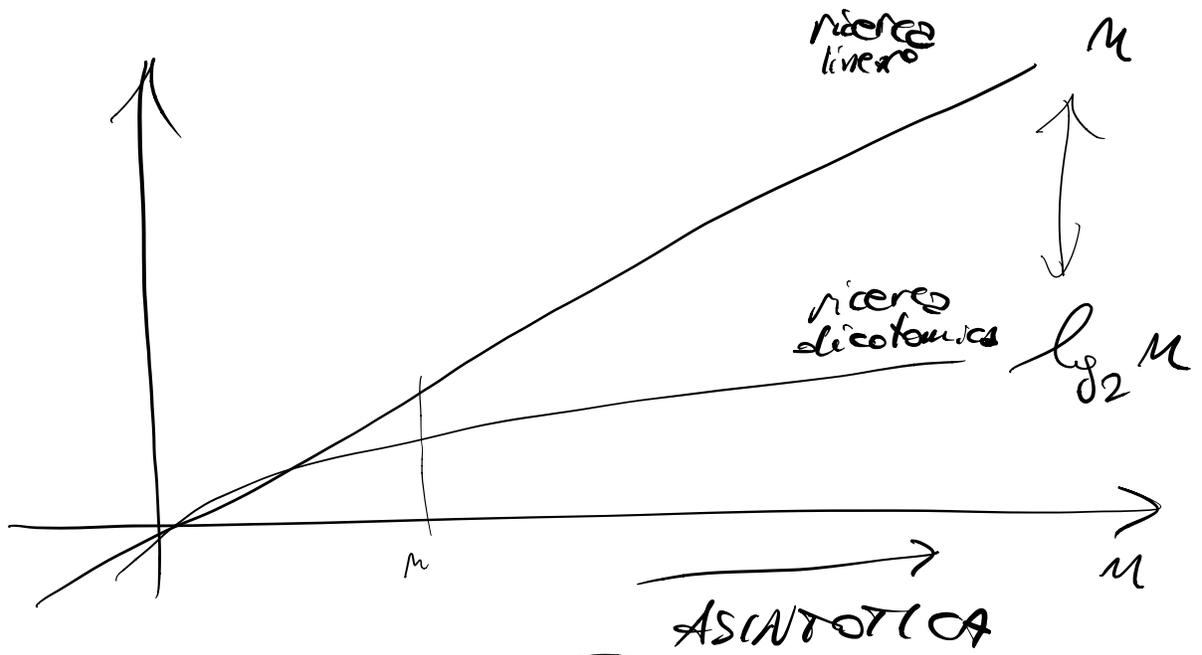
$$n < 2^k$$

$$\log_2 n < \log_2 2^k = k \log_2 2 = k$$

$$\log_2 n < k$$

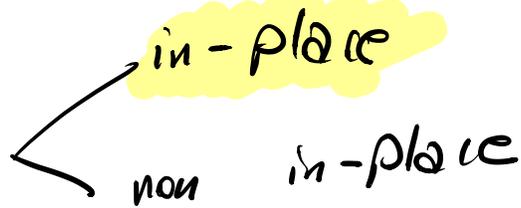
$$\lceil \log_2 n \rceil$$

~~Caso precedente  $\rightarrow \lceil \log_2 n \rceil$~~   
~~CONFRONTI~~

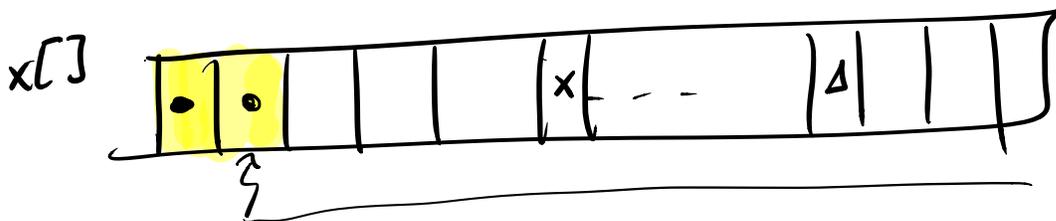


$$\lim_{n \rightarrow \infty} \frac{\log n}{n} = 0$$

# ORDINAMENTO



## SELECTION SORT

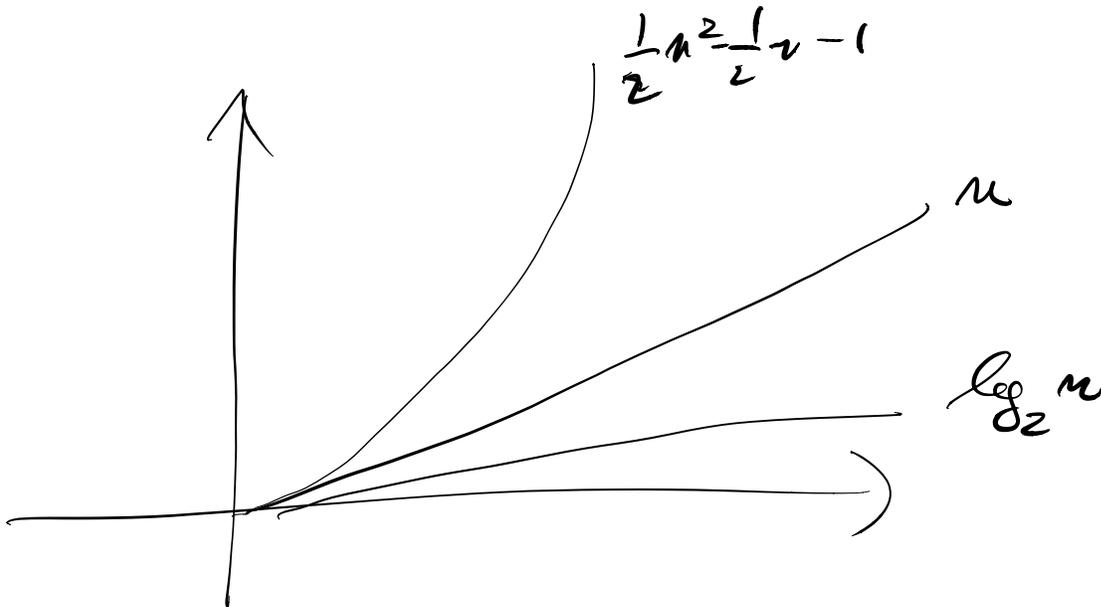


```
func selectionSort(x [int]) {  
  // Posizione da sistemare  
  for i=0; i < n; i++ {  
    imin=i  
    min = x[i]  
    for j=i+1; j < n; j++ {  
      if x[j] < min {  
        imin=j  
        min = x[j]  
      }  
    }  
    x[i], x[imin] = x[imin], x[i]  
  }  
}
```

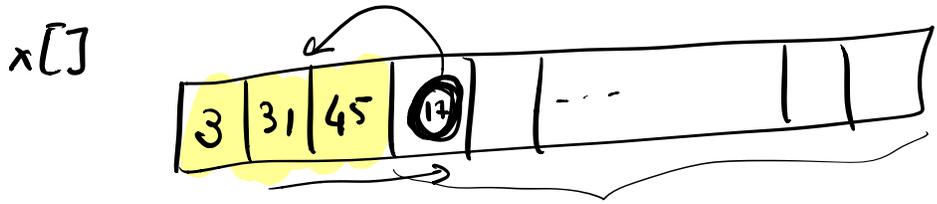
$$\begin{aligned}
 & n-1 \quad n-2 \quad n-3 \quad \dots \quad 2 \\
 2 + 3 + \dots + (n-1) &= \sum_{i=1}^{n-1} i - 1 = \\
 &= \frac{(n-1)n}{2} - 1 = \frac{1}{2}n^2 - \frac{1}{2}n - 1
 \end{aligned}$$

CASO PEGGIORE: n° confronti:

$$\frac{1}{2}n^2 - \frac{1}{2}n - 1 = O(n^2)$$



# INSERTION SORT



```

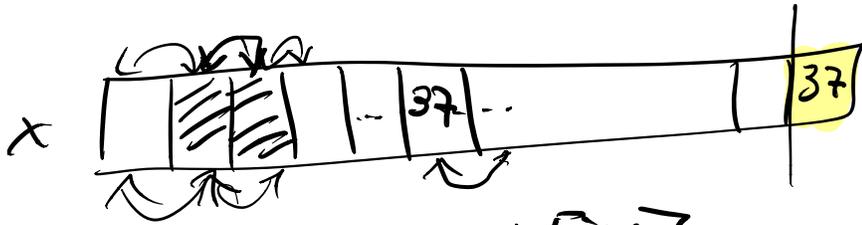
func insertionSort(x [int]) {
    var n = len(x)
    for i := 1; i < n; i++ {
        j := i+1 // Do insertion from i+1
        for k := 0; k < i; k++ {
            if x[k] > x[i] {
                break
            }
        }
        temp := x[j]
        for h := i; h >= k; h-- {
            x[h+1] = x[h]
        }
        x[k] = temp
    }
}
    
```

$$\begin{array}{ccccccc}
 1 & 2 & 3 & \dots & n-1 & & \\
 \underbrace{\hspace{10em}} & & & & & & \\
 1+2+3+\dots+n-1 & = & \sum_{i=1}^{n-1} i & = & \frac{(n-1)(n)}{2} & = & \\
 & & & & & & \\
 & & & & = \frac{1}{2}n^2 - \frac{1}{2}n & & 
 \end{array}$$

Also PEGGcore:  $n^{\circ}$  assignments:

$$\frac{1}{2}n^2 - \frac{1}{2}n$$

# BUBBLE SORT



$$\rightarrow x[i] \leq x[i+1]$$

```

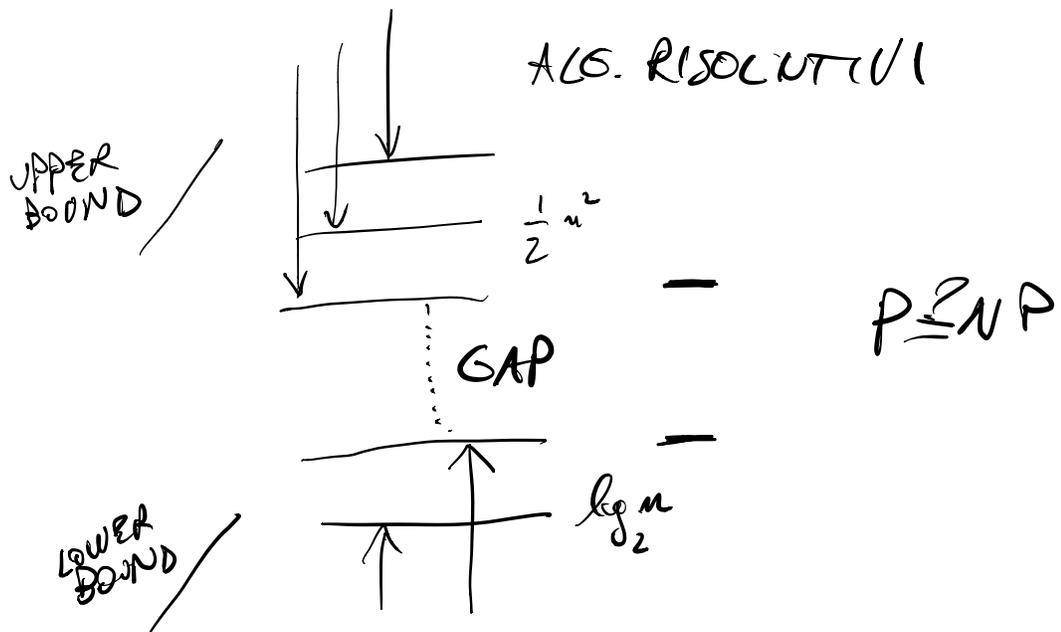
func bubbleSort (x [int]) {
  for (j = n; j > 0; j--) {
    for (i = 0; i < j-1; i++) {
      if (x[i] > x[i+1]) {
        x[i], x[i+1] = x[i+1], x[i]
      }
    }
  }
}
  
```

$$\begin{aligned}
 & (n-1) + (n-2) + \dots + 1 + 0 = \\
 & = \sum_{i=1}^{n-1} i = \frac{(n-1)n}{2} = \frac{1}{2}n^2 - \frac{1}{2}n
 \end{aligned}$$

CASO PEGGIORE

$$\frac{1}{2} n^2 - \frac{1}{2} n$$

LOWER/UPPER BOUND



- ORDINAMENTO CON SCAMBI  
E CONFRONTI

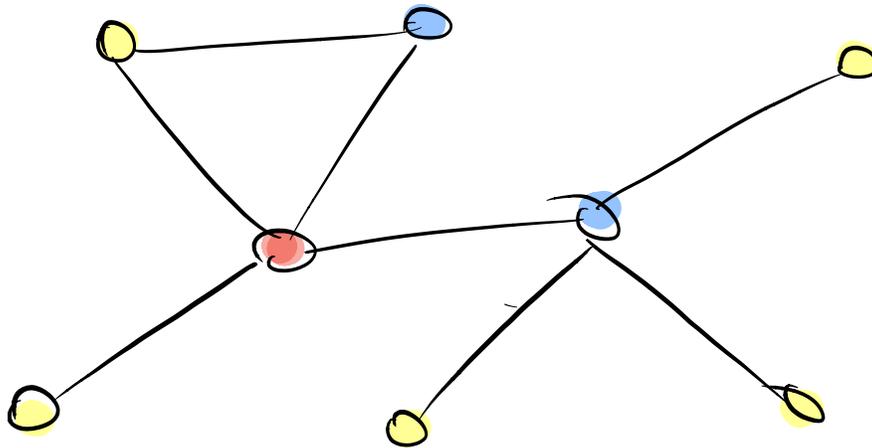
RICHIEDE  $\Omega(n \lg n)$

- ESISTONO ALGORITMI  $O(n \lg n)$

{ - HEAPSORT  
- QUICKSORT

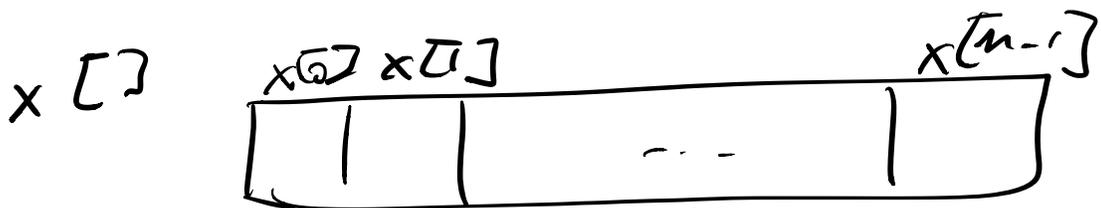
# ALGORITMO NAIF $O(n^2)$

COLORABILITY ~ NP-completi



$$2^M$$

$$M$$



```

type persona struct {
    nome, cognome string
    data Nascita date
    reddito Anno int
}
var x []persona
    
```

